

Conceptual modelling, ontology design, and semantic interoperability

Professional master on technologies for e-government

Lecture 7

Claudio Masolo

Laboratory for Applied Ontology, ISTC-CNR

`masolo@loa-cnr.it`

Trento, 22.10.2007

Formal ontology

Brings together two senses of *formal*.

- ▷ **Rigorous:** formal framework, e.g., logic.
- ▷ **Foundational:** search for invariants across domains (Husserl).

Rigorous

Starting from the informal conceptual analysis it is necessary to choose:

- ▷ the kinds of entities we accept in the domain;
- ▷ the formal language (logical framework + primitives);
- ▷ the properties of the primitives (axioms).

Foundational

Basic tools for ontology development:

- ▷ import results coming from *logic, philosophy, linguistics, cognitive science, mathematics, computer science*, etc. in order to identify a set of well founded conceptual primitives that are general enough to be applied to several domains constituting the basis for the development of more domain oriented models.
- ⇒ Intrinsically multidisciplinary.

Formal ontology again

Facilitates:

- ▷ a precise representation (and communication) based on cognitively transparent and theoretically well founded primitives;
- ▷ the **reuse** of theories in different contexts and domains;
- ▷ a modular approach to ontology building;
- ▷ the **comparison** between different theories that correspond to alternative ontological positions.

Basic notions

But what are these basic notions coming from different disciplines?

- ▷ Essentiality and identity.
- ▷ Properties and qualities.
- ▷ Dependence.
- ▷ Parthood, connection, and unity.
- ▷ Time, space, and space-time.
- ▷ Objects, constitution, and composition.
- ▷ Events, participation, and causality.

Outline

- ▷ Orders
- ▷ Mereologies
- ▷ Theories of time

Orders are quite simple and abstract structures that are useful for general and heterogeneous comparisons between entities:

- ▷ “is more/less ... than”: “is bigger than”, “is smaller than”, “is later than” ...
- ▷ “is to the left of”, “is an ancestor of”, “is a divisor of”, “is part of” ...

Note: in the following we will consider as formal framework the FOL with identity.

▷ A **partial order** is a binary relation \leq such that:

$$\mathbf{A1} \quad x \leq x \quad (\textit{reflexivity})$$

$$\mathbf{A2} \quad x \leq y \wedge y \leq z \rightarrow x \leq z \quad (\textit{transitivity})$$

$$\mathbf{A3} \quad x \leq y \wedge y \leq x \rightarrow x = y \quad (\textit{antisymmetry})$$

▷ A **strict partial order** is a binary relation $<$ such that:

$$\mathbf{A4} \quad x < y \wedge y < z \rightarrow x < z \quad (\textit{transitivity})$$

$$\mathbf{A5} \quad x < y \rightarrow \neg y < x \quad (\textit{asymmetry})$$

$$\mathbf{T1} \quad \neg x < x \quad (\textit{irreflexivity})$$

▷ In a partial order, defining $<$ as:

$$\mathbf{D1} \quad x < y \triangleq x \leq y \wedge \neg y \leq x$$

we obtain a strict partial order.

▷ In a strict partial order, defining \leq as

$$\mathbf{D2} \quad x \leq y \triangleq x < y \vee x = y$$

we obtain a partial order.

▷ Therefore it is the same to start from the \leq or the $<$ primitive.

▷ Classical additional properties for orders

A6 $x \leq y \vee y \leq x$ *(total/linear order)*

A7 $x < y \rightarrow \exists z(x < z \wedge z < y)$ *(dense order)*

A8 $x < y \rightarrow \exists z(x < z \leq y \wedge \neg \exists v(x < v < z))$
 $x < y \rightarrow \exists z(x \leq z < y \wedge \neg \exists v(z < v < y))$ *(discrete order)*

A9 $\exists x \forall y(x \leq y)$ *(left bounded)*

A10 $\exists x \forall y(y \leq x)$ *(right bounded)*

A11 $\forall x \exists y(y < x)$ *(left unbounded)*

A12 $\forall x \exists y(x < y)$ *(right unbounded)*

▷ On finite domains, all orders are discrete and bounded.

▷ **Examples:** $\langle \mathbb{N}, \leq \rangle$, $\langle \mathbb{Z}, \leq \rangle$, $\langle \mathbb{Q}, \leq \rangle$, $\langle \mathbb{R}, \leq \rangle$, $\langle 2^A, \subseteq \rangle$.

- ▷ Many orders on many domains!
 - weights, heights, numbers, instants (precedence), preferences... these are not parthood!
- ▷ Specificities of Parthood?
 - Surely not a linear order
 - Dense, discrete, bounded, unbounded: all possible options

So?

A bit of history of mereologies

Mereology: from the greek meros, ‘the theory of parthood’

- ▷ Lesniewski 1927-1931, On the Foundations of Mathematics. Alternative to Set Theory for escaping Russells paradox
 - No null individual (no empty set)
 - No distinction between *urelements* (\in) and *sets* (\subseteq): a single relation of parthood
- ▷ Tarski 1935. Link with algebra.
- ▷ Leonard and Goodman 1940. The calculus of individuals, nominalism.
- ▷ Contemporary studies: Peter Simons (1986), Achille Varzi (1996).

Why is the Parthood relation important?

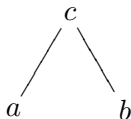
Today, the majority of ontologies use at least a parthood relation.

- ▷ Philosophical, cognitive and linguistic relevance.
- ▷ Spatial and temporal reasoning based on vague information: impossibility to use exact coordinates, trajectories in terms of mathematical functions, and calculus.
- ▷ Reference to extended entities (e.g., temporal periods, spatial regions), possibly composed of parts of the same nature.
- ▷ No calculus, yet still a rigorous formal approach: logical theories.

Note: similarly to the case of orders, no one single mereology, but a plurality of different mereologies.

Hasse diagrams

- ▷ Graph on finite domains
- ▷ **Convention:** all vertical or oblique arcs are implicitly oriented from bottom to top and indicate a strict order.
- ▷ **Example:** the following diagram graphically represents the formula $a < c \wedge b < c$



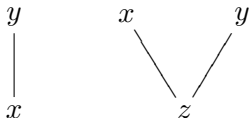
Basic mereology (M)

- ▷ We note the parthood relation with P .
- ▷ A *basic mereology* M is a reflexive, transitive, and antisymmetric relation, i.e. it is a partial order.
- ▷ Useful definitions (*proper part* corresponds to strict order):

$$\mathbf{D3} \quad PP(x, y) \triangleq P(x, y) \wedge \neg P(y, x) \quad (\textit{proper part})$$

$$\mathbf{D4} \quad O(x, y) \triangleq \exists z(P(z, x) \wedge P(z, y)) \quad (\textit{overlap})$$

$$\mathbf{D5} \quad PO(x, y) \triangleq O(x, y) \wedge \neg P(x, y) \wedge \neg P(y, x) \quad (\textit{proper overlap})$$

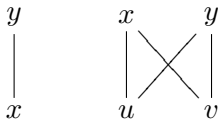


Minimal and extensional mereologies (1/2)

▷ Supplementation

$$\text{ws} \quad \text{PP}(x, y) \rightarrow \exists z(\text{P}(z, y) \wedge \neg \text{O}(z, x)) \quad (\text{weak suppl.})$$

$$\text{ss} \quad \neg \text{P}(x, y) \rightarrow \exists z(\text{P}(z, y) \wedge \neg \text{O}(z, x)) \quad (\text{strong suppl.})$$



▷ Extensionality:

$$\text{we} \quad \exists z(\text{PP}(z, x)) \wedge \forall z(\text{PP}(z, x) \leftrightarrow \text{PP}(z, y)) \rightarrow x = y \quad (\text{weak})$$

$$\text{se} \quad \forall z(\text{O}(z, x) \leftrightarrow \text{O}(z, y)) \rightarrow x = y \quad (\text{strong})$$

note that **we** does not follow from **ws**, see theorem (T5).

Minimal and extensional mereologies (2/2)

$$\text{MM} = \text{M} \cup \{\text{ws}\}$$

minimal mereology

$$\text{EM} = \text{M} \cup \{\text{ss}\}$$

extensional mereology

▷ Theorems:

$$\mathbf{T2} \quad \text{EM} \vdash \text{ws}$$

$$\mathbf{T3} \quad \text{EM} \vdash \text{we}$$

$$\mathbf{T4} \quad \text{EM} \vdash \text{se}$$

$$\mathbf{T5} \quad \text{MM} \not\vdash \text{we}$$

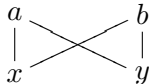
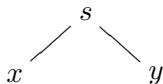
$$\mathbf{T6} \quad \text{M} \cup \{\text{we}\} \not\vdash \text{se}$$

$$\mathbf{T7} \quad \text{M} \cup \{\text{se}\} \vdash \text{ws}$$

▷ Sum:

$$\mathbf{D6} \quad \text{Sum}(s, x, y) \triangleq \forall z (\text{O}(z, s) \leftrightarrow (\text{O}(z, x) \vee \text{O}(z, y)))$$

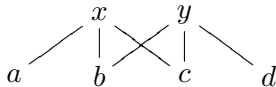
$$\text{sum} \quad \forall x, y \exists s (\text{Sum}(s, x, y))$$



▷ Product:

$$\mathbf{D7} \quad \text{Prod}(p, x, y) \triangleq \forall z (\text{P}(z, p) \leftrightarrow (\text{P}(z, x) \wedge \text{P}(z, y)))$$

$$\text{prd} \quad \text{O}(x, y) \rightarrow \exists p (\text{Prod}(p, x, y))$$



▷ Difference:

$$\begin{aligned} \mathbf{D8} \quad \text{Diff}(d, x, y) &\triangleq \forall z (P(z, d) \leftrightarrow (P(z, x) \wedge \neg O(z, y))) \\ \text{dif} \quad \exists z (P(z, x) \wedge \neg O(z, y)) &\rightarrow \exists d (\text{Diff}(d, x, y)) \end{aligned}$$

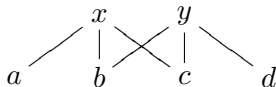
▷ Complement:

$$\begin{aligned} \mathbf{D9} \quad \text{Compl}(c, x) &\triangleq \forall z (P(z, c) \leftrightarrow \neg O(z, x)) \\ \text{cmp} \quad \exists z (\neg O(z, x)) &\rightarrow \exists c (\text{Compl}(c, x)) \end{aligned}$$

▷ Universe (noted with Un):

$$\text{uni} \quad \exists x \forall y (P(y, x))$$

$$\mathbf{T8} \quad \exists z (z = \text{Un}) \wedge x \neq \text{Un} \rightarrow (\text{Compl}(c, x) \leftrightarrow \text{Diff}(c, \text{Un}, x))$$



Closure mereology

(3/3)

$$\text{CM} = \text{M} \cup \{\text{sum}, \text{prd}\}$$

closure mereology

$$\text{CMM} = \text{CM} \cup \{\text{ws}\}$$

closure minimal mereology

$$\text{CEM} = \text{CM} \cup \{\text{ss}\}$$

closure extensional mereology

▷ Theorems:

T9 $\text{CMM} \vdash \text{ss}$ therefore CMM and CEM are equivalent.

T10 $\text{CEM} \vdash \text{Sum}(s, x, y) \wedge \text{Sum}(s', x, y) \rightarrow s = s'$

T11 $\text{CEM} \vdash \text{Prod}(p, x, y) \wedge \text{Prod}(p', x, y) \rightarrow p = p'$

T12 $\text{CEM} \vdash \text{dif}$

T13 $\text{CEM} \vdash \text{Diff}(d, x, y) \wedge \text{Diff}(d', x, y) \rightarrow d = d'$

T14 $\text{CEM} \vdash \text{Compl}(c, x) \wedge \text{Compl}(c', x) \rightarrow c = c'$

Atomic mereology

▷ Atoms, atomicity, and atomic essentialism:

D10 $\text{At}(x) \triangleq \forall y(\text{P}(y, x) \rightarrow y = x)$ (*atom*)

at $\forall x \exists y(\text{At}(y) \wedge \text{P}(y, x))$ (*atomicity*)

ate $(\text{At}(z) \rightarrow (\text{P}(z, x) \rightarrow \text{P}(z, y))) \rightarrow \text{P}(x, y)$ (*at. essent.*)

▷ Atomic mereologies

$\text{AEM} = \text{EM} \cup \{\text{at}\}$

the same for all other mereologies

▷ Theorems:

T15 $\text{AEM} \vdash \text{ate}$

T16 $\text{EM} \cup \{\text{ate}\} \vdash \text{at}$

T17 $\text{M} \cup \{\text{ate}\} \vdash \text{ss}$

T18 $\text{AM} \cup \{\text{we}\} \vdash \text{se}$

Back to mereology requirements

- ▷ No distinction between *urelements* (\in) and *sets* (\subseteq)
 - in the Sum “operator”, the *sum* and the *addenda* are at the same level
- ▷ No null individual (no empty set)
 - the mereologies we considered are not based on an axiom like

$$\mathbf{A13} \quad \exists x(\forall y(P(x, y)))$$

Which mereology?

▷ **Extensionality**

- Identity between my body and the collection of my organs
- parthood vs. spatial inclusion + theory of constitution and levels of reality

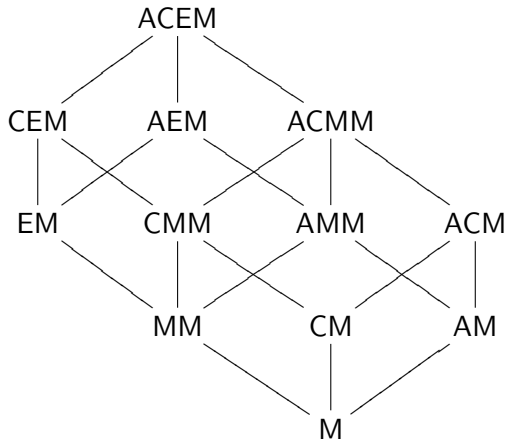
▷ **Closure**

- sum of my nose and Caesars toe
- mereotopology (part+connection) to identify *wholes*

▷ **Transitivity**

- My hand is part of me, and
I'm part of the University of Trento, but
My hand is not part of the Committee

The space of mereologies



Libraries of theories

- ▷ The *libraries* of theories as, for example, the library/space of mereologies are very important *conceptual* tools because they ‘encapsulate’ a deep analysis of different notions that are intuitively linked to the a general one (parthood for mereologies).
- ▷ Different theories can be *adequate* to specific modeling requirements: the user selects the theory that better matches his needs.
- ▷ No monolithic/standardized approach: the links between the theories in the library make explicit their (in)compatibilities.
 - In this lecture, we will see how different mereologies can be *reused* for modeling some perspectives on time.
 - In the following of the course you will see their application to: space, physical objects, qualities, organizations, etc.