THE BDI MODEL OF AGENCY AND BDI LOGICS

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Introduction – Beliefs, Desires, Intentions (Bratman)

- Desires and beliefs range over states of affairs, while intentions range over actions and by extension, plans.
- Intentions are persistent, whereas desires can be dropped at any time.
- Intentions need not be holded forever.
- Intentions drive means-end reasoning.
- Beliefs constrain desires.
- Intentions constrain future deliberation and planning.
- Intentions influence beliefs upon which future practical reason is based.
- Intentions imply a degree of commitment to a goal.
- Intentions and beliefs are required to be consistent, i.e. not to imply some kind of pragmatical contradiction. This condition is assumed to imply that of rationality.
- Intentions, beliefs and desires need not be complete or, to put it simply, all-encompassing.
- Beliefs are subject to revision.
- Intentions and hence plans can be reconsidered.

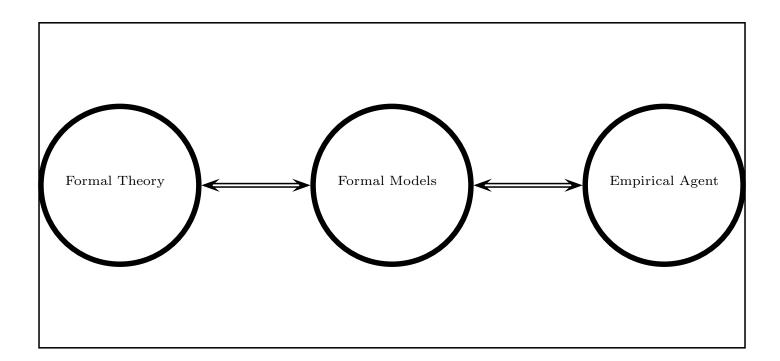
Introduction – Agents (Bratman)

Definition 1 We understand by an <u>agent</u> an entity (a moral or a legal person, a computer program) that is capable of reacting to a certain environment through its performing a certain number of actions over which it can exert some kind of control. We say further that an agent is <u>rational</u> if his actions, decisions, plans and intentions are consistent or coherent with his beliefs and desires as well as between themselves.

Introduction – Kinds of Agents (Bratman)

- *Blindly-minded agents* are agents that are blindy over-committed to their basic beliefs and intentions or desires, which they never put in question nor revise. They can be seen as fanatical agents.
- Single-minded agents are agents whose (derived) intentions may change due to belief revision. They are thus cautiously committed to their intentions. They can be seen as cautious agents. They are able to modify a plan if needed.
- *Open-minded agents* are agents that revise their beliefs and that change both their desires and derived intentions accordingly. They are thus under-committed to their intentions.

Introduction – Modelling Desiderata



BDI Logic Language Primitives

Definition 2 The set \mathcal{F} of the formulae of BDI logic is defined by the grammar:

- $< o var > ::= x_1 | ... | x_n, n \ge 0.$
- $< e var > ::= e_1 | ... | e_m, m \ge 0.$
- $< pred > ::= P_0^0 |...| P_k^l, \, k, l \ge 0.$
- $\langle var \rangle ::= \langle o var \rangle | \langle e var \rangle.$
- < atom > ::= < pred > (< var >, ..., < var >).
- < state form > ::= succeeded(< event var >)|failed(< event var >)| $| <math>< atom > |\neg < state - form > | < state - form > \lor < state - form > |$ $| <math>\exists < var > < state - form > |$ **Bel** < state - form > || $\mathbf{Go} < state - form > |$ **In** < state - form > |optional < path - form >.
- $< path form > ::= < state form > |\neg < path form > | < path form > \lor < path form > | < path form > U < path form > | < path form > U < path form > | < path form > .$
- < form > ::= < state form > | < path form >.

Formal Semantics – Models

Definition 3 A model for BDI logics is a Kripke branching-time temporal model with three distinct accessibility relations. Let \mathcal{R} denote the set of propositional symbols. Then a model is a structure $M = (D, E, T, W; \prec, I, B, G; \Phi)$ where:

- D is a non-empty set called domain of objects.
- E is a non-empty set called domain of events.
- T is a non-empty set of time points.
- $\prec \subseteq T \times T$ is the branching time relation.
- W is a non empty set of worlds over T.
- $I \subseteq W \times T \times W$ an intention accessibility relation.
- $B \subseteq W \times T \times W$ a belief accessibility relation.
- $G \subseteq W \times T \times W$ a goal accessibility relation.
- $\Phi: \mathcal{R} \times W \times T \to \bigcup_{i \in \mathbb{N}} \wp(D^i)$ is an interpretation function for predicate symbols.

Formal Semantics – Satisfaction

Definition 4 The satisfaction relation is defined by induction on \mathcal{F} as follows – first on path formulas and then on state formulas. Let v be an assignation and v^* be an assignation identical to v but for some object or event variable x or e. Then:

- $M \models_{w_t}^v P(x_1, ..., x_n)$ iff $(v(x_1), ..., v(x_n)) \in \Phi(P, w, t)$.
- $M \models_{w_t}^v \neg A \text{ iff } M \nvDash_{w_t}^v A.$
- $M \models_{w_t}^{v} A \lor B$ iff $M \models_{w_t}^{v} A$ or $M \models_{w_t}^{v} B$.
- $M \models_{w_t}^v \exists x^O A \text{ iff } M \models_{w_t}^{v^*} A \text{ for some } d \in D_O.$
- $M \models_{w_t}^v \exists x^E A \text{ iff } M \models_{w_t}^{v^*} A \text{ for some } e \in D_E$
- $M \models^{v}_{\langle w_{t_0}, w_{t_1}, \ldots \rangle} A \text{ iff } M \models^{v}_{w_{t_0}} A.$
- $M \models^{v}_{< w_{t_0}, w_{t_1}, \ldots >} \bigcirc A \text{ iff } M \models^{v}_{< w_{t_1}, \ldots, >} A$
- $M \models_{< w_{t_0}, w_{t_1}, \ldots >}^{v} \Diamond A$ iff for some $i \ge 0$ such that $M \models_{< w_{t_i}, \ldots, >}^{v} A$.
- $M \models_{<w_{t_0},w_{t_1},\ldots>}^{v} AUB$ iff either of these conditions hold:
 - 1. For some $i \ge 0$ such that $M \models_{< w_{t_i}, \ldots >}^{v} B$ and for all $0 \le j < i, M \models_{< w_{t_i}, \ldots >}^{v} A.$
 - 2. For any $j \ge 0, M \models_{\langle w_{t_j}, \dots \rangle}^{v} A$.
- $M \models_{w_{t_0}}^v optional A \text{ iff for some fullpath} < w_{t_0}, w_{t_1}, \ldots >,$ $M \models_{<w_{t_0}, w_{t_1}, \ldots >}^v A.$

- $M \models_{w_t}^v succeed(e)$ iff for some time point t', $S_w(t', t) = v(e)$.
- $M \models_{w_t}^v failed(e)$ iff for some time point t', $F_w(t', t) = v(e)$.
- $M \models_{w_t}^{v} \mathbf{Bel}A$ iff for any $w' \in \mathcal{B}_t^w$, $M \models_{w'_t}^{v} A$.
- $M \models_{w_t}^v \mathbf{In}A$ iff for any $w' \in \mathcal{I}_t^w, M \models_{w'_t}^v A$.
- $M \models_{w_t}^{v} \mathbf{Go}A$ iff for any $w' \in \mathcal{G}_t^w$, $M \models_{w'_t}^{v} A$.

Example of Model

The following model M_1 satisfyies the formula

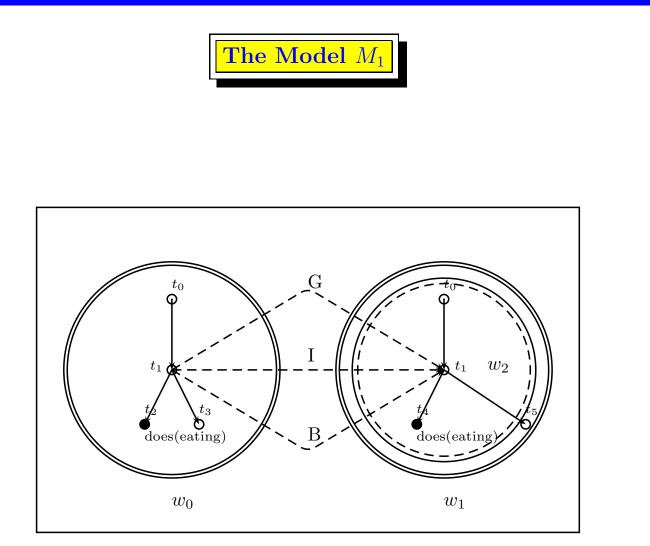
$$\mathbf{In}(optional(\diamondsuit(does(eating)))) \to \mathbf{Go}(optional(\diamondsuit(does(eating))))$$

For indeed we have that:

 $M_1 \models_{w_{0_{t_1}}} \mathbf{In}(optional(\diamondsuit(does(eating)))))$

implies that:

 $M_1 \models_{w_{0_{t_1}}} \mathbf{Go}(optional(\diamondsuit(does(eating)))).$



 M_1

Correspondence Theory

Name	Modal Formula Scheme	Condition on Model
BK	$\mathbf{Bel}(A \to B) \to (\mathbf{Bel}A \to \mathbf{Bel}B)$	B non empty.
BD	$\mathbf{Bel}A \to \neg \mathbf{Bel} \neg A$	B is serial.
B4	$\mathbf{Bel}A \to \mathbf{Bel}\mathbf{Bel}A$	B is transitive.
B5	$\neg \mathbf{Bel} \neg A \to \mathbf{Bel} \neg \mathbf{Bel} \neg A$	B is euclidian.
IK	$\mathbf{In}(A \to B) \to (\mathbf{In}A \to \mathbf{In}B)$	I non empty.
ID	$\mathbf{In}A \to \neg \mathbf{In} \neg A$	I is serial.
GK	$\mathbf{Go}(A \to B) \to (\mathbf{Go}A \to \mathbf{Go}B)$	G non empty.
GD	$\mathbf{Go}A \to \neg \mathbf{Go} \neg A.$	G is serial.
G-B	$\mathbf{Go}A \to \mathbf{Bel}A$	$G \subseteq B.$
I-G	$\mathbf{In}A ightarrow \mathbf{Go}A$	$I \subseteq G.$
G-B*	$\mathbf{Go} \alpha \to \mathbf{Bel} \alpha$	$G \subseteq_{struct} B.$
I-G*	$\mathbf{In} \alpha ightarrow \mathbf{Go} \alpha$	$I \subseteq_{struct} G.$

Basic Axioms

Name	Modal Formula Scheme	Intuitive property
BK	$\mathbf{Bel}(A \to B) \to (\mathbf{Bel}A \to \mathbf{Bel}B)$	Belief implication closure
BD	$\mathbf{Bel}A \to \neg \mathbf{Bel} \neg A$	Belief consistency
B4	$\mathbf{Bel}A \to \mathbf{BelBel}A$	Belief positive introspection
B5	$\neg \mathbf{Bel} \neg A \to \mathbf{Bel} \neg \mathbf{Bel} \neg A$	Belief negative introspection
IK	$\mathbf{In}(A \to B) \to (\mathbf{In}A \to \mathbf{In}B)$	Intention implication closure
ID	$\mathbf{In}A \to \neg \mathbf{In} \neg A$	Intention consistency
GK	$\mathbf{Go}(A \to B) \to (\mathbf{Go}A \to \mathbf{Go}B)$	Goal implication closure
GD	$\mathbf{Go}A \to \neg \mathbf{Go} \neg A.$	Goal consistency
G-B*	$\mathbf{Go}\alpha \to \mathbf{Bel}\alpha$	Desire-belief compatibility
I-G*	$\mathbf{In}lpha ightarrow \mathbf{Go}lpha$	Intention-desire compatibility

Intention Axioms

Name	Axiom Scheme	Intuitive property
A11	$\mathbf{In}A \to \mathbf{Bel}(\mathbf{In}A)$	Intentions about beliefs
A12	$\mathbf{Go}A \to \mathbf{Bel}(\mathbf{Go}A)$	Beliefs about goals
A13	$\mathbf{In}A ightarrow \mathbf{Go}(\mathbf{In}A)$	Desires about intentions
A14	$\forall e(\mathbf{In}(does(e)) \to does(e)$	Intentions leading to ac-
		tions
A15	$\forall e(done(e)) \rightarrow \mathbf{Bel}(done(e))$	Awareness of primitive
		events
A16	$\mathbf{In}A \to inevitable(\diamondsuit(\neg \mathbf{In}A))$	No infinite deferral prop-
		erty

Committeent Axioms

Name	Axiom Scheme	Intuitive property
C1	$\mathbf{In}(inevitable(\diamondsuit A)) \rightarrow$	Blind-mindedness
	$inevitable(\mathbf{In}(inevitable(\diamondsuit A))\mathbf{U}A)$	property
C2	$\mathbf{In}(inevitable(\diamondsuit A)) \rightarrow$	Single-mindedness
	$inevitable(\mathbf{In}(inevitable(\diamondsuit A))$	property
	$\mathbf{U}(A \vee \neg \mathbf{Bel}(optional(\diamondsuit A)))$	
C3	$\mathbf{In}(inevitable(\diamondsuit A)) \rightarrow$	Open-mindedness
	$inevitable(\mathbf{In}(inevitable(\diamondsuit A))$	property
	$\mathbf{U}(A \vee \neg \mathbf{Go}(optional(\diamondsuit A)))$	

The Systems

Definition 5 The axioms without the commitment axioms constitute the I system: the <u>Basic</u> Intention System. They model basic agents. Commitment axioms extend it to cover different kinds of agents.



- $I \vdash \mathbf{In}\alpha \to \mathbf{Bel}\alpha$
- $I \vdash (\mathbf{Go}\alpha \wedge \mathbf{In}(\alpha \to \beta)) \to \mathbf{Go}\beta.$
- $I \vdash (\mathbf{Bel}\alpha \wedge \mathbf{In}(\alpha \to \beta)) \to \mathbf{Bel}\beta.$
- $I \vdash (\mathbf{Bel}\alpha \wedge \mathbf{Go}(\alpha \to \beta)) \to \mathbf{Bel}\beta.$

Extensions

Name	Modal Formula Scheme	Intuitive property
E1	$\neg(\mathbf{In}A \wedge \mathbf{Bel} \neg A)$	Intention-belief consistency
E2	$\mathbf{In}A\wedge \neg BelA$	Intention-belief incompleteness
E3	$\mathbf{Bel}A\wedge \neg \mathbf{Go}A$	Transference property
E4	$\mathbf{In}A \wedge \mathbf{Bel}(A \to B)$	Side effects property
	$\wedge \neg \mathbf{In}A$	



Some extensions do not work, for example:

 $I + E2 \vdash \bot$

- (i.e. Bratman's asymmetry thesis) Since:
- 1. $InA \rightarrow BelA$ Proposition 2.3
- 2. \neg (**Go** $A \land \neg$ **Bel**A) 1,PL
- 3. $InA \land \neg BelA E2$
- 4. $\perp -2,3, PL$



- There are only two primitive BDI modalities: **GOAL** (desire) and **BEL** (belief). The intention modality, i.e. **INT**, is defined by imposing persistence conditions on **GOAL**.
- The language includes variables (and constants) ranging over agents.
- There are action modalities together with action connectives. Control structures (iterative and conditional) can be defined.
- Possible worlds in models are discrete linear orders (finite or inifite) with a least lower bound. Furthermore, the accessiblity relations satisfy strong realism.
- No formal system is given.



- We ignore if full BDI logics are sound or complete, although we assume that they are sound. They are, anyway, undecidable.
- Different axioms convey different properties of agents.
- Formal constrainment of future desires, beliefs and intentions by present intentions is not very intuitive.
- They are a good specification tool (used to develop expremental software agents).
- Not all of Bratman's *provisi* and thesis hold.
- BDI logics model only deliberation and not means-end reasoning.

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